

**Question A – Recursion and Induction – 20 Marks**

Given the sequence  $a_n$  defined with the recurrence relation:

$$a_0 = 1$$

$$a_n = n (a_{n-1})^2 \text{ for } n \geq 1$$

Terms of a Sequence (5 marks)

$$a_1 = 1 \cdot 1^2 = 1$$

$$a_2 = 2 \cdot 1^2 = 2$$

$$a_3 = 3 \cdot (2 \cdot 1^2)^2 = 3 \cdot 2^2 \cdot 1^4 = 12$$

$$a_4 = 4 \cdot (3 \cdot 2^2 \cdot 1^4)^2 = 4 \cdot 3^2 \cdot 2^4 \cdot 1^8 = 576$$

$$a_5 = 5 \cdot (4 \cdot 3^2 \cdot 2^4 \cdot 1^8)^2 = 5 \cdot 4^2 \cdot 3^4 \cdot 2^8 \cdot 1^{16} = 1,658,880$$

Iteration (3 marks)

$$a_n = \prod_{i=1}^n i^{2^{n-i}} = \prod_{i=0}^{n-1} (n-i)^{2^i}$$

Proof by induction (12 marks)

Show that  $2 \mid n^2 - n$  for all positive integers  $n$  by weak induction. No other method is acceptable.

Define the predicate  $P(n)$  to be  $2 \mid n^2 - n$ .

We are going to show that  $\forall n \in \mathbb{N}^+ P(n)$

Proof:

Base case:

When  $n=1$ ,  $n^2 - n = 1 - 1 = 0 = 2 \cdot 0$ , so  $2 \mid n^2 - n$

i.e.  $P(1)$  is true.

Inductive Step:

Assume that  $P(n)$  is true for some  $n$  in  $\mathbb{N}^+$

This means that  $\exists k \in \mathbb{Z}$ ,  $n^2 - n = 2k$

Show that  $P(n+1)$  is true.

$(n+1)^2 - (n+1) = n^2 + 2n + 1 - n - 1 = n^2 - n + 2n$  by algebra

$= 2k + 2n$  by inductive hypothesis

$= 2(k+n)$  by algebra

Since  $k, n$  are integers and  $\mathbb{Z}$  is closed under  $+$ , then  $k+n \in \mathbb{Z}$

Therefore  $2 \mid (n+1)^2 - (n+1)$

QED by induction

**Question B – Number Theory – 20 marks**Euclidian Algorithm (5 marks)

$$\begin{aligned} \gcd(598, 1287) &= \gcd(1287, 598) \\ &= \gcd(598, 1287 \bmod 598) = \gcd(598, 91) \\ &= \gcd(91, 598 \bmod 91) = \gcd(91, 52) \\ &= \gcd(52, 91 \bmod 52) = \gcd(52, 39) \\ &= \gcd(39, 52 \bmod 39) = \gcd(39, 13) \\ &= \gcd(13, 39 \bmod 13) = \gcd(13, 0) = 13 \end{aligned}$$

Mod Proof (15 marks)

Prove that for any integers  $A, B, a, d$  such that  $d \neq 0$ ,  
if  $A \bmod d = a$  and  $B \bmod d = 1$  then  $AB \bmod d = a$

Proof:

Let  $A, B, a, d$  be integers such that  $d \neq 0$  and  $A \bmod d = a$  and  $B \bmod d = 1$

We'll show that  $AB \bmod d = a$

Since  $A \bmod d = a$  then by the QRT:  $A = (A \text{ div } d) \cdot d + a$  (1) and  $0 \leq a < d$  (2)

Since  $B \bmod d = 1$  then by the QRT:  $B = (B \text{ div } d) \cdot d + 1$  (3)

So  $AB = ((A \text{ div } d) \cdot d + a) \cdot ((B \text{ div } d) \cdot d + 1)$  by substituting with (1) and (3)  
 $= [(A \text{ div } d)(B \text{ div } d) \cdot d + (B \text{ div } d) \cdot a + (A \text{ div } d)] \cdot d + a$

Let  $p = [(A \text{ div } d)(B \text{ div } d) \cdot d + (B \text{ div } d) \cdot a + (A \text{ div } d)]$

Then  $AB = p \cdot d + a$  (4)

Since all the terms in  $p$  are integers and  $\mathbb{Z}$  is closed under  $\text{div}$ ,  $+$  and  $\cdot$ , then  $p \in \mathbb{Z}$ . (5)

By the QRT:  $\exists!q, r \in \mathbb{Z} \ AB = q \cdot d + r$  and  $0 \leq r = AB \bmod d < d$

Therefore, since by (4):  $AB = p \cdot d + a$  and by (2):  $0 \leq a < d$ , where  $p, a \in \mathbb{Z}$  by (5)

then  $a$  must be  $AB \bmod d$  (because of the uniqueness part of the QRT)

QED

**Question C – Graph Theory – 20 marks**

Graph Degrees (12 marks)

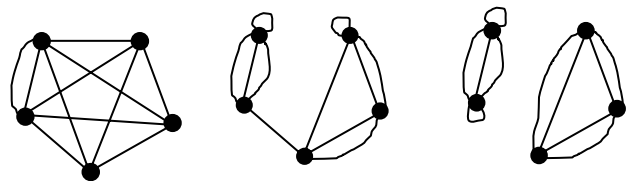
For each of the following questions, either draw a graph with the requested properties, or explain **convincingly** (possibly by quoting a theorem) why such a graph cannot be drawn.

a) A graph with 5 vertices of degrees 5, 5, 4, 4, 3

This graph cannot be drawn because the degree of a graph (sum of degree of vertices) must be even, and in this case the sum of the degree of these 5 vertices is 21 which is odd.

b) A graph with 5 vertices of degrees 5, 5, 4, 4, 4

There are many possible answers. Here are 3:



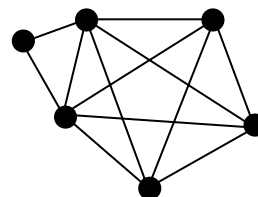
c) A simple graph with 5 vertices of degrees 5, 5, 4, 4, 4

Simple graphs do not contain parallel edges or loops, so the maximum degree of any edge in a simple graph with  $n$  vertices is  $n-1$  as each vertex has at most  $n-1$  incident edges, one for each of the other vertices.

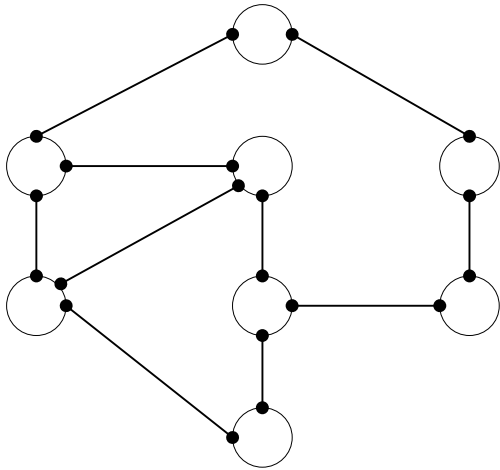
Therefore it is not possible to draw a simple graph with  $n$  vertices, at least one of which has degree  $n$ ,

d) A simple graph with at least 5 vertices which have degrees 5, 5, 4, 4, 4. The other vertices have whichever degree seems appropriate.

Again there are many possible answers. One simple answer is built by adding one vertex to  $K_5$  and connecting it to 2 of  $K_5$ 's vertices, thus bringing their degrees up from 4 to 5:



Circuits (8 marks)



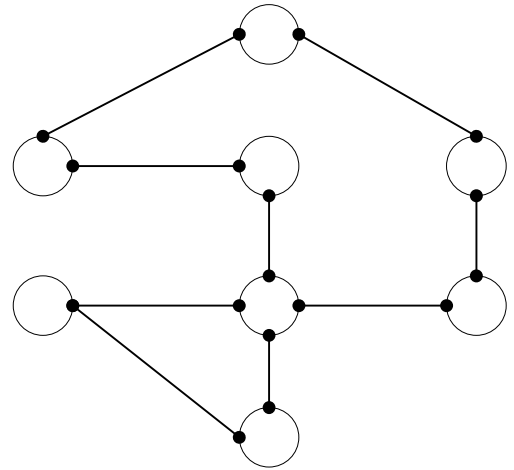
a) Find an Euler circuit in G that starts at  $v_1$ .

A graph contains an Euler circuit iff it is connected and every vertex has an even degree. G, however, has some vertices with odd degree ( $v_2, v_3, v_5, v_6$ ). Therefore it cannot contain an Euler circuit.

b) Find a Hamiltonian circuit in G that starts at  $v_1$

There are 2 Hamiltonian circuits:

- $v_1 e_1 v_2 e_3 v_3 e_5 v_5 e_{10} v_8 e_8 v_6 e_9 v_7 e_4 e_2 v_1$
- $v_1 e_2 v_4 e_7 v_9 e_6 e_8 v_8 e_{10} v_5 e_5 v_3 e_3 v_2 e_1 v_1$



c) Find an Euler circuit in H that starts at  $v_1$

There are 4 Euler circuits:

- $v_1 e_2 v_4 e_3 v_7 e_7 v_6 e_6 v_5 e_8 v_8 e_9 v_6 e_4 v_3 e_3 v_2 e_1 v_1$
- $v_1 e_2 v_4 e_3 v_7 e_7 v_6 e_9 v_8 e_8 v_5 e_6 v_6 e_4 v_3 e_3 v_2 e_1 v_1$
- $v_1 e_1 v_2 e_3 v_3 e_4 v_6 e_6 v_5 e_8 v_8 e_9 v_6 e_7 v_7 e_5 v_4 e_2 v_1$
- $v_1 e_1 v_2 e_3 v_3 e_4 v_6 e_9 v_8 e_8 v_5 e_6 v_6 e_7 v_7 e_5 v_4 e_2 v_1$

d) Find a Hamiltonian circuit in H that starts at  $v_1$

Theorem: If H contained a Hamiltonian circuit, then it would have a connected subgraph which contained all its vertices, all of which would be of degree 2.

To get such a subgraph, all the edges incident on vertices in  $V(H) - \{v_6\}$  would have to be preserved because all the vertices in that set have degree 2. So the subgraph would have to include edges  $e_4, e_6, e_7, e_9$  and as a result vertex  $v_6$  would have to be of degree 4 which contradicts the theorem.

v1

e1

v3

e3

e4

e5

v5

v6

e10

v8